Meaning 5.2: Semantics 5 — Toy models; applications of what we've learned

May 20, 2020

Recap

Model architecture: based on set theory.

Basic units: elements/individuals, ordered pairs, sets

TRUTH CONDITIONS

 $\frac{\mathsf{Operations/relationships:}}{\mathsf{subset} \subseteq, \; \mathsf{superset} \supseteq} \; \; \mathsf{membership} \in , \; \mathsf{intersection} \; \cap, \; \mathsf{union} \; \cup, \\ \mathsf{COMPOSITIONALITY}$

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Denotations we've covered so far

The following linguistic units are represented by these objects in our model:

- (i) names, definite NPs (DPs) elements/individuals a,b,c,\ldots ${
 m Ted}{
 m I}=t$
- (ii) nouns, adjectives, intransitive verbs sets of individuals A, B, C, ... $[dance] = \{x: x \text{ dances}\}$
- (iii) transitive verbs sets of ordered pairs $\mathcal{A},\mathcal{B},\mathcal{C},\ldots$ $[like] = \{\langle x,y \rangle: \ y \ likes \ x\}$

Compositions

- (i) simple predication S \rightarrow DP VP; VP \rightarrow V or VP = is + AdjP S is true iff $\llbracket DP \rrbracket \in \llbracket VP \rrbracket$
- $\begin{tabular}{ll} \begin{tabular}{ll} \b$
- (iii) definite descriptors $\mathsf{DP} \to \mathsf{D} \; \mathsf{NP}$
 - $[\![\mathsf{D}\ \mathsf{NP}]\!] = \mathsf{contextually}\ \mathsf{salient}\ \mathsf{individual}\ d\ \mathsf{s.t.}\ d \in [\![\mathsf{NP}]\!]$

Quantifiers

The quantification determiner used (all, some, no,...) determines the relevant relationship to check for the sets denoted by $[\![N]\!]$ and $[\![VP]\!]$

There is usually some domain restriction (determined by the context)

Simple predication is not used in this case of S \rightarrow DP VP

- (i) All N VP is true iff $[\![N]\!] \subseteq [\![VP]\!]$, such that $[\![N]\!]$ consists of the contextually salient individuals that have that property.
- (ii) Some N VP is true iff $([N] \cap [VP]) \neq \emptyset$
- (iii) No N VP is true iff ($[\![N]\!] \cap [\![VP]\!]$) = \varnothing

Checking against a model

To see if our semantic theory is working, we can build a toy model to check our theory.

Our model has individuals, sets of individuals and sets of ordered pairs as its units, and the operations of \cup , \cap and the relations \in , \subseteq .

We can build a hypothetical world w, and specify the individuals and sets in the world w, to check if our theories are producing the right predictions.

The idea here is that we don't **have** to precisely model the world at large to know something about language: all we need to do is model a context where we would know the behavior of language, and that is a much easier task. So, we typically just specify a 'toy model' to play around with and observe behavior. Ideally, the behavior of our model will reflect linguistic behavior.

A toy model

In the model world M, let's assume there are five individuals:

And some of the predicates are as follows:

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\begin{array}{c} \text{philosopher: } \{a,f,p,r\} \\ \text{linguist: } \{c\} \\ \text{tall:} \{a,c,f\} \\ \text{know:} \{\langle a,p\rangle, \langle f,r\rangle, \langle r,f\rangle, \langle p,a\rangle, \langle r,c\rangle\} \end{array}
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Statements to check against this model

Let's start rattling off things we get from this model.

• Frege is a philosopher.

Russell is tall.

All linguists are philosophers.

Statements to check against this model

The linguist is a tall.

Plato knows is Aristotle.

Some philosophers are tall.

 $\begin{tabular}{l} [Some philosophers are tall] = true iff [philosopher] \cap [tall] $\neq \varnothing$ \\ \hline TRUE \\ \end{tabular}$

Some interesting cases

• The philosopher is a tall.

[The philosopher is tall] = true iff [The philosopher] \in [tall] TRUE?, FALSE?

only defined if there is contextually salient individual in the set of philosophers ... presupposition.

Everybody knows somebody.

Everybody knows somebody = true iff (i) every person knows at least one person; (ii) there is at least one person that every other person knows.

TRUE on (i), FALSE on (ii)

scope ambiguity with quantifiers every and some; see colors on next slides.

Everybody knows somebody — reading (i), true

(i) Everybody knows somebody.

Everybody knows somebody = true iff every person knows at least one person

TRUE

Why? — For all y in the complete set of individuals in the world w, there is at least one x such that there is a pair $\langle x,y\rangle$ such that $\langle x,y\rangle\in \text{know}:\{\langle a,p\rangle,\langle f,r\rangle,\langle r,f\rangle,\langle p,a\rangle,\langle r,c\rangle\}$

<u>Intuition</u>: are all the individuals in world w represented on the righthand side of the pairs $\langle left, right \rangle$ in $[\![know]\!]$?

Everybody knows somebody — reading (ii), false

(ii) Everybody knows somebody.

Everybody knows somebody = true iff there is at least one person that every other person knows.

FALSE

Why? — There is not one x such that for all y in the complete set of individuals in the world w, it is the case that there exists a pair $\langle x,y\rangle$ such that $\langle x,y\rangle\in \text{know}:\{\langle a,p\rangle,\langle f,r\rangle,\langle r,f\rangle,\langle p,a\rangle,\langle r,c\rangle\}$

<u>Intuition</u>: Is there a lefthand element such that the pair $\langle left, right \rangle$ exists for all individuals y in the model world M when you substitute that individual in for the righthand element?

End of this video's material. End of semantics and my lecturing. Thank you *so* much for your attention and effort. Please stay safe and healthy — good luck and hope to see you around.

Cheers,

Brandon